

Jumping Abstraction Layers: Microarchitectural Attacks in JavaScript

Daniel Gruss

September 18, 2019

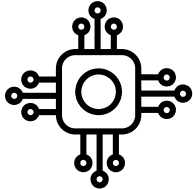
Graz University of Technology





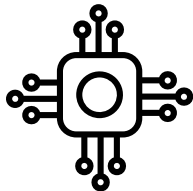
side channel
= obtaining meta-data and
deriving secrets from it

CHANGE MY MIND



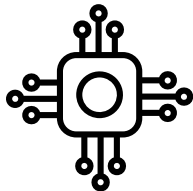
Microarchitecture...

- not architectural state



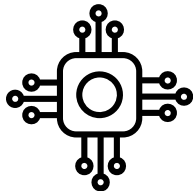
Microarchitecture...

- not architectural state
- not visible to software



Microarchitecture...

- not architectural state
- not visible to software
- hardware specific



Microarchitecture...

- not architectural state
- not visible to software
- hardware specific
- changes with generations









1337 4242

FOOD CACHE

Revolutionary concept!

Store your food at home,
never go to the grocery store
during cooking.

Can store **ALL** kinds of food.

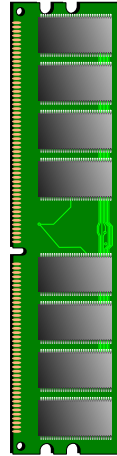
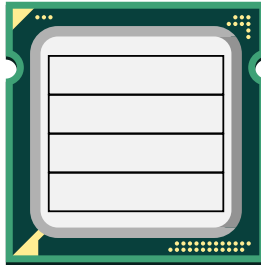
ONLY TODAY INSTEAD OF ~~\$1,300~~

\$1,299

ORDER VIA PHONE: +555 12345

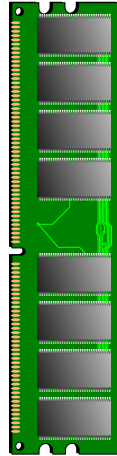
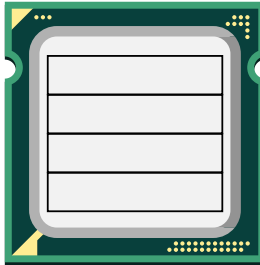


```
printf("%d", i);  
printf("%d", i);
```



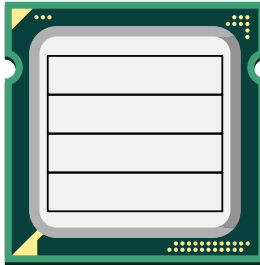
```
printf("%d", i);  
printf("%d", i);
```

Cache miss

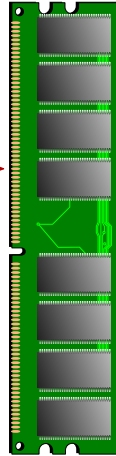


```
printf("%d", i);  
printf("%d", i);
```

Cache miss

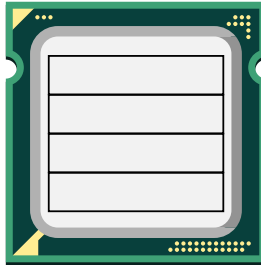


Request



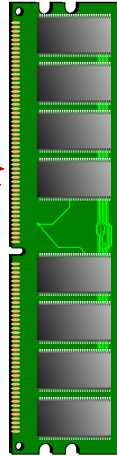
```
printf("%d", i);  
printf("%d", i);
```

Cache miss



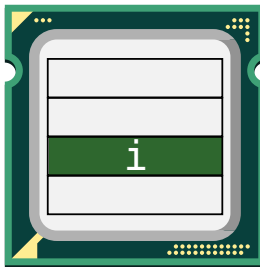
Request

Response



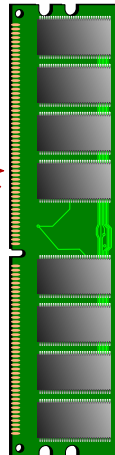

```
printf("%d", i);  
printf("%d", i);
```

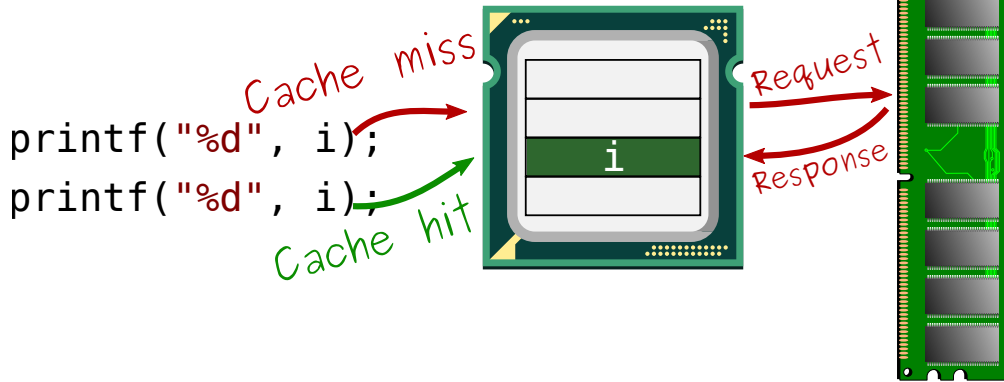
Cache miss



Request

Response



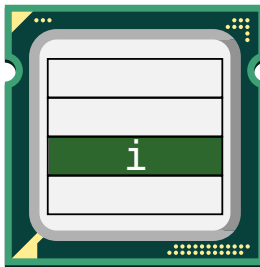


DRAM access,
slow

```
printf("%d", i);  
printf("%d", i);
```

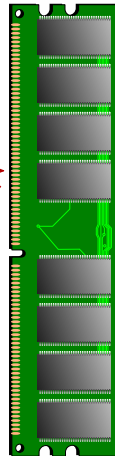
Cache miss

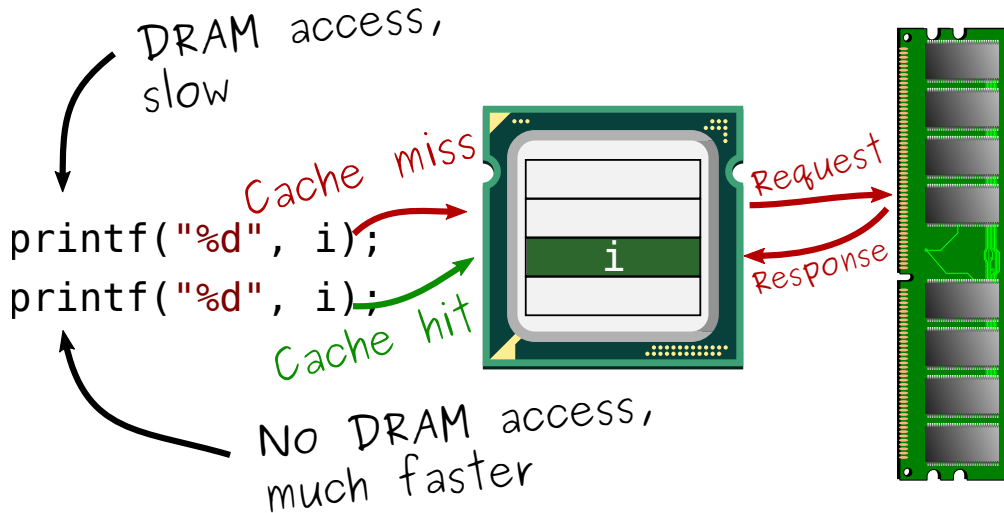
Cache hit



Request

Response

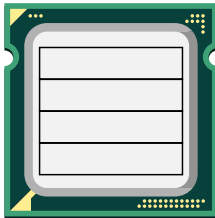




Shared Memory

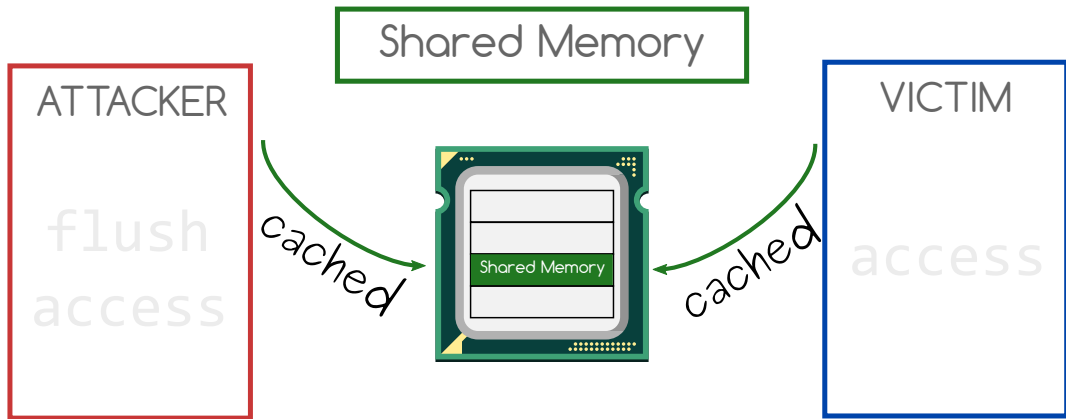
ATTACKER

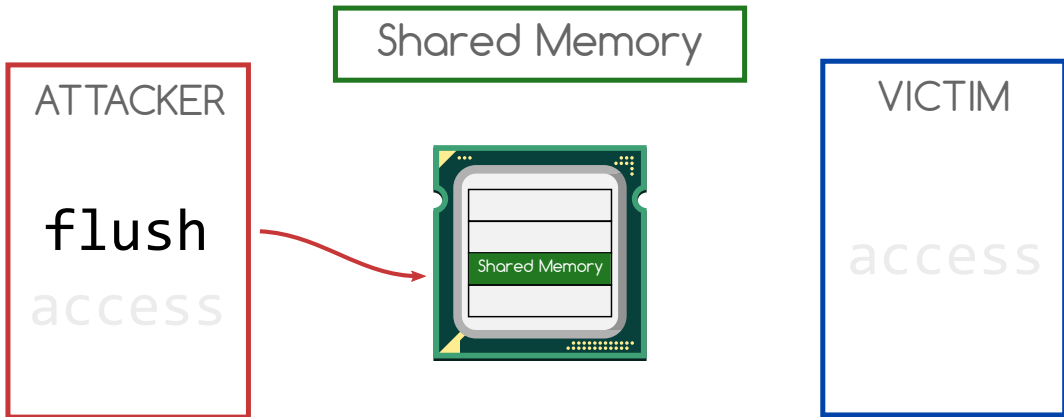
flush
access

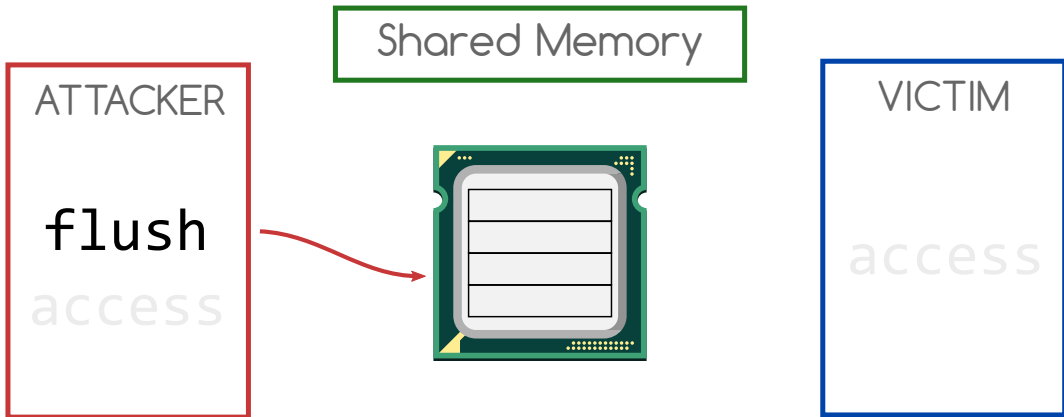


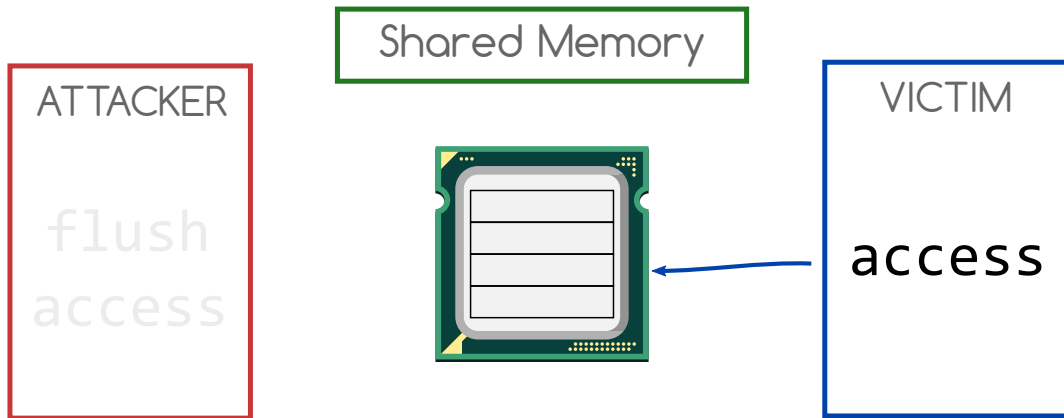
VICTIM

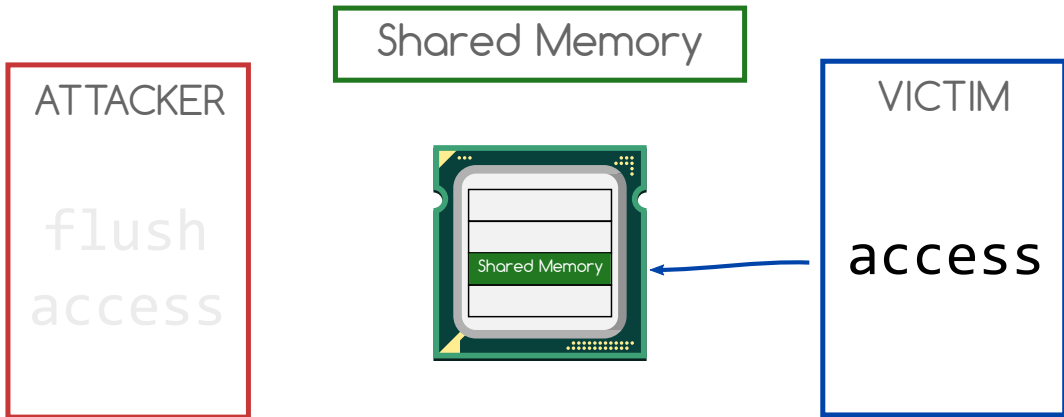
access

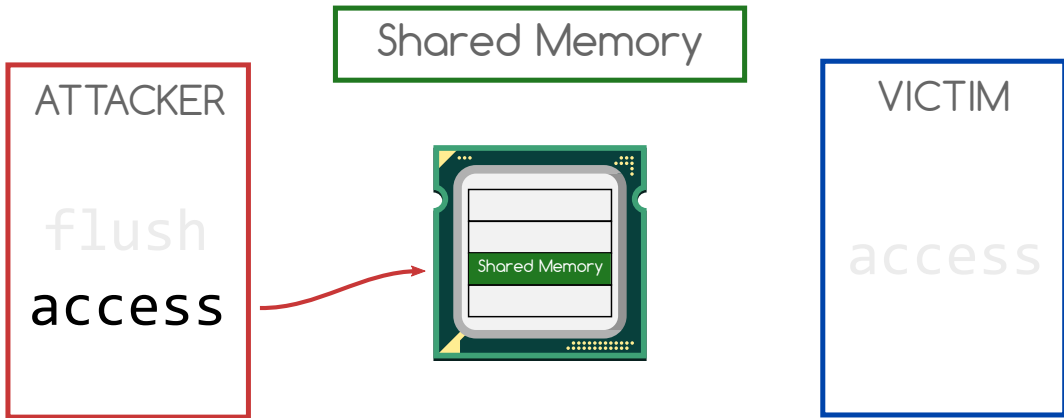


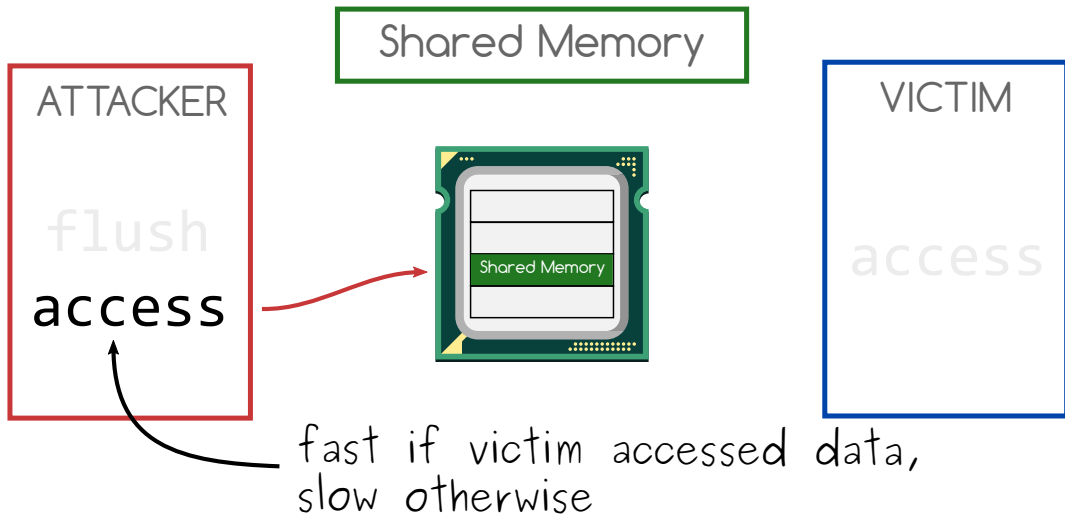






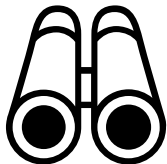






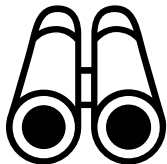
TIMING...

IT'S ALL ABOUT



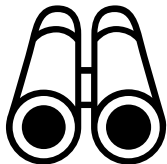
First **side-channel** attack in JavaScript

- Stone et al. (2013): HTML5 pixel perfect attacks



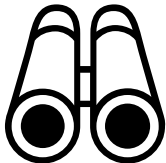
First **side-channel** attack in JavaScript

- Stone et al. (2013): HTML5 pixel perfect attacks
- Use high-resolution timer



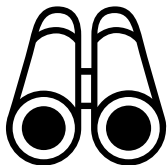
First **side-channel** attack in JavaScript

- Stone et al. (2013): HTML5 pixel perfect attacks
- Use high-resolution timer
- Timing redraw events (visited, ...)



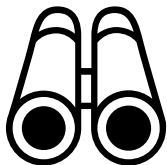
First **side-channel** attack in JavaScript

- Stone et al. (2013): HTML5 pixel perfect attacks
- Use high-resolution timer
- Timing redraw events (visited, ...)
- SVG filter timing for pixels (known since 2011)



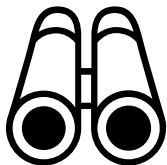
First **microarchitectural** attack in JavaScript

- Oren et al. (2015): The Spy in the Sandbox



First **microarchitectural** attack in JavaScript

- Oren et al. (2015): The Spy in the Sandbox
- Timing of memory accesses



First **microarchitectural** attack in JavaScript

- Oren et al. (2015): The Spy in the Sandbox
- Timing of memory accesses
- Data cached or not

FANTASTIC TIMERS

AND WHERE
TO FIND THEM

HIGH-RESOLUTION MICROARCHITECTURAL
ATTACKS IN JAVASCRIPT



- We need a high-resolution timer

- We need a **high-resolution timer**
- Native: `rdtsc`

- We need a **high-resolution timer**
- Native: `rdtsc`
- JavaScript: `performance.now()`


- We need a **high-resolution timer**
- Native: `rdtsc`
- JavaScript: `performance.now()`

`performance.now()`

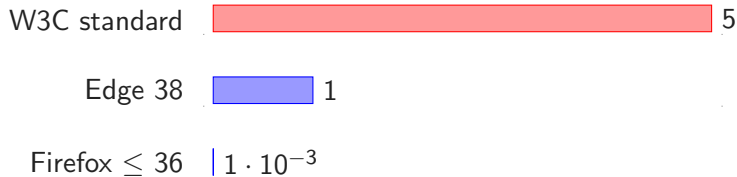
[...] represent times as floating-point numbers with up to microsecond precision.

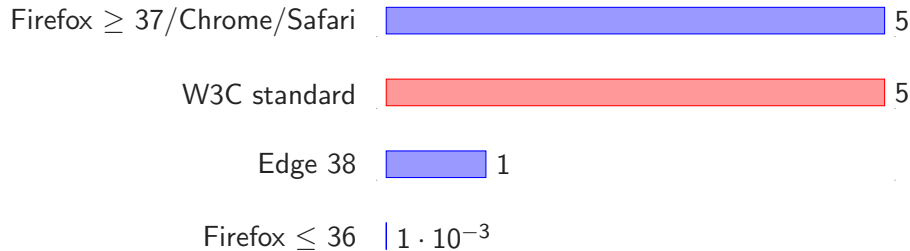
— Mozilla Developer Network

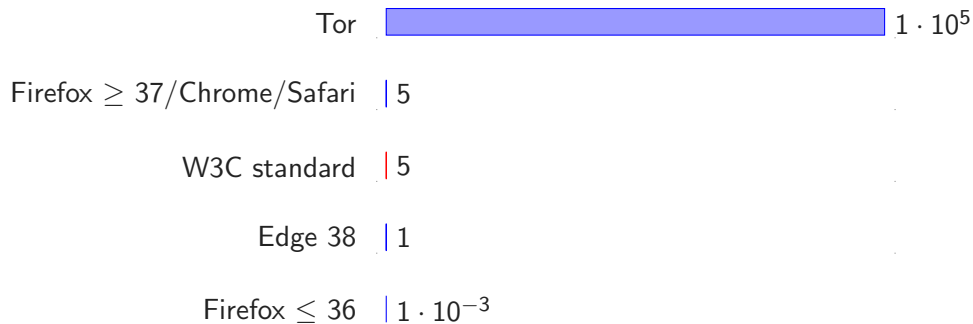
Firefox $\leq 36 \cdot 10^{-3}$

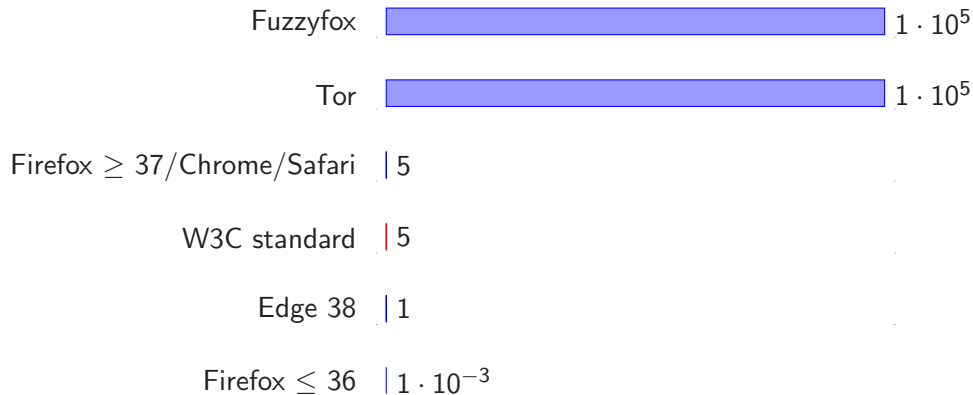
Edge 38  1

Firefox ≤ 36 | $1 \cdot 10^{-3}$











- Current precision can't measure **cycle differences**



- Current precision can't measure **cycle differences**
- Two options



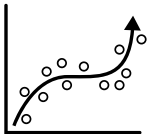
- Current precision can't measure **cycle differences**
- Two options
- **Recover** a higher resolution

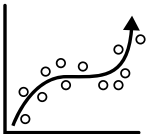


- Current precision can't measure **cycle differences**
- Two options
- **Recover** a higher resolution
- **Build** our own high-resolution timer

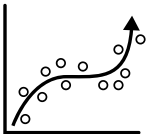


- **Measure** how often we can **increment** a variable between two timer ticks

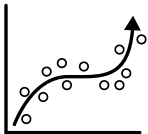




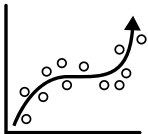
- **Measure** how often we can **increment** a variable between two timer ticks
- Average number of increments is the **interpolation step**



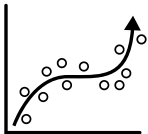
- **Measure** how often we can **increment** a variable between two timer ticks
- Average number of increments is the **interpolation step**
- To measure with high resolution:



- **Measure** how often we can **increment** a variable between two timer ticks
- Average number of increments is the **interpolation step**
- To measure with high resolution:
 - Start measurement at **clock edge**



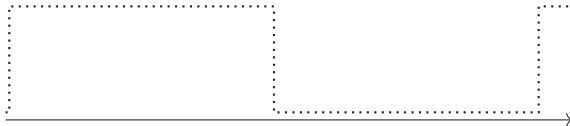
- **Measure** how often we can **increment** a variable between two timer ticks
- Average number of increments is the **interpolation step**
- To measure with high resolution:
 - Start measurement at **clock edge**
 - **Increment** a variable until next clock edge



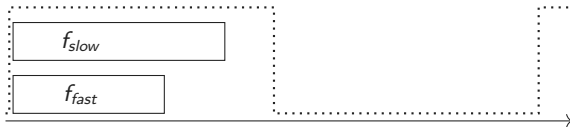
- **Measure** how often we can **increment** a variable between two timer ticks
- Average number of increments is the **interpolation step**
- To measure with high resolution:
 - Start measurement at **clock edge**
 - **Increment** a variable until next clock edge
- Highly accurate: 500 ns (Firefox/Chrome), 15 μ s (Tor)

- We can get a higher resolution for a **classifier** only

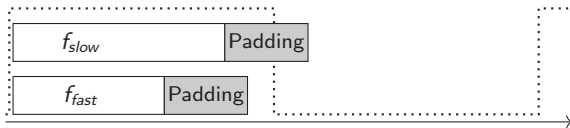
- We can get a higher resolution for a **classifier** only
- Often sufficient to see which of two functions takes **longer**



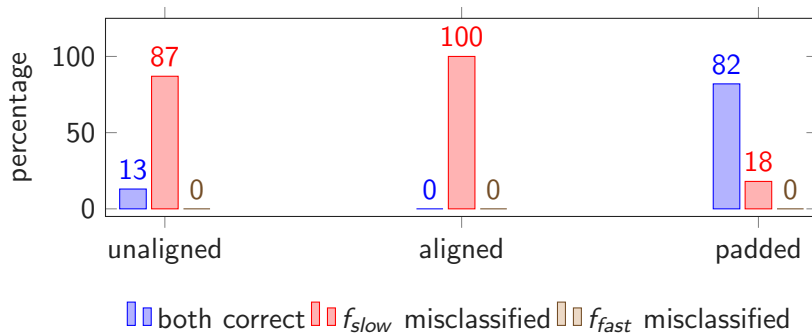
- We can get a higher resolution for a **classifier** only
- Often sufficient to see which of two functions takes **longer**

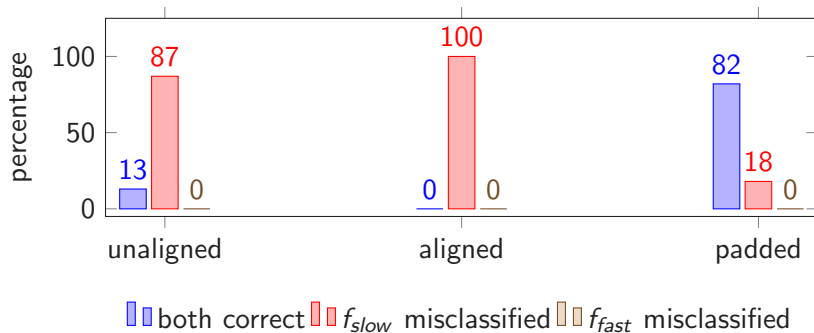


- We can get a higher resolution for a **classifier** only
- Often sufficient to see which of two functions takes **longer**

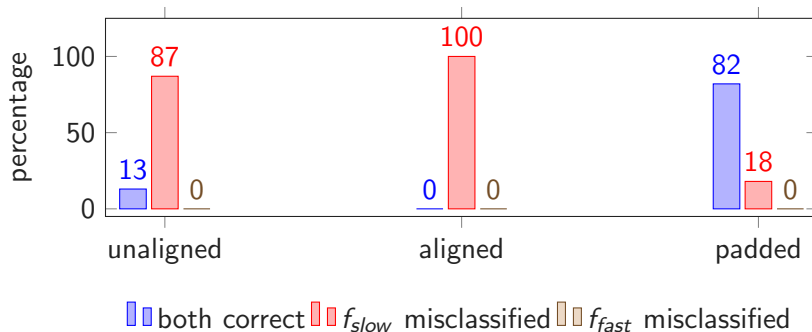


- **Edge thresholding**: apply padding such that the slow function crosses one more clock edge than the fast function.

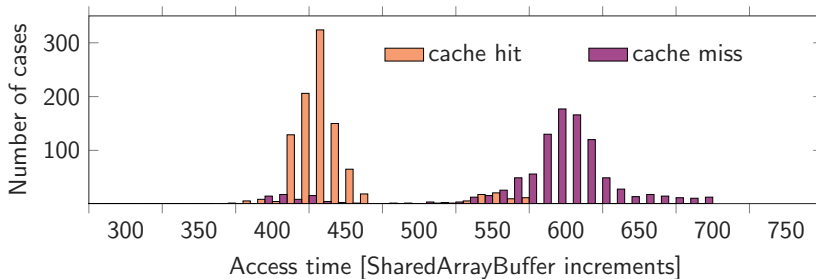


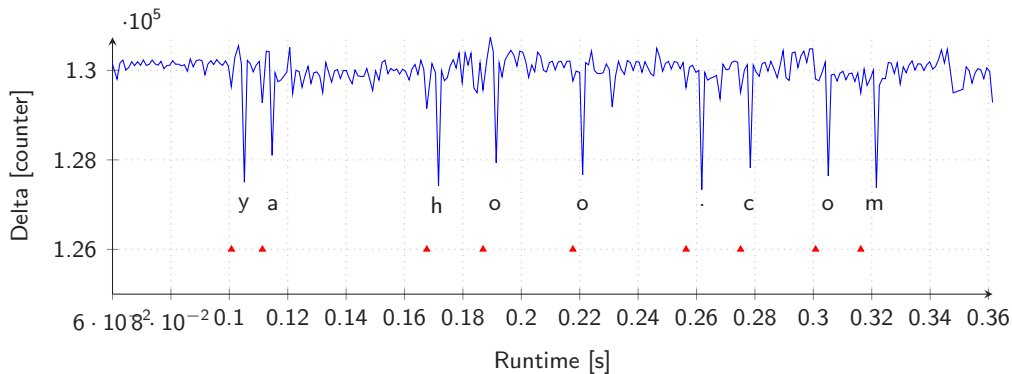


- Yields **nanosecond** resolution



- Yields **nanosecond** resolution
- Firefox/Tor (2 ns), Edge (10 ns), Chrome (15 ns)





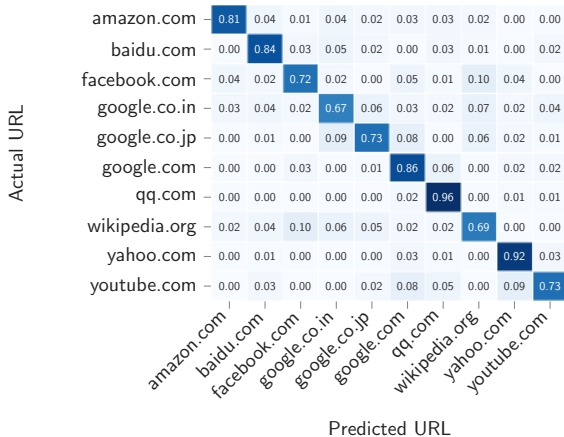


Figure 1: Confusion matrix for URL input.

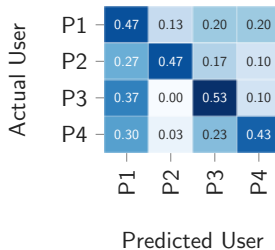


Figure 2: Confusion matrix for input by different users.

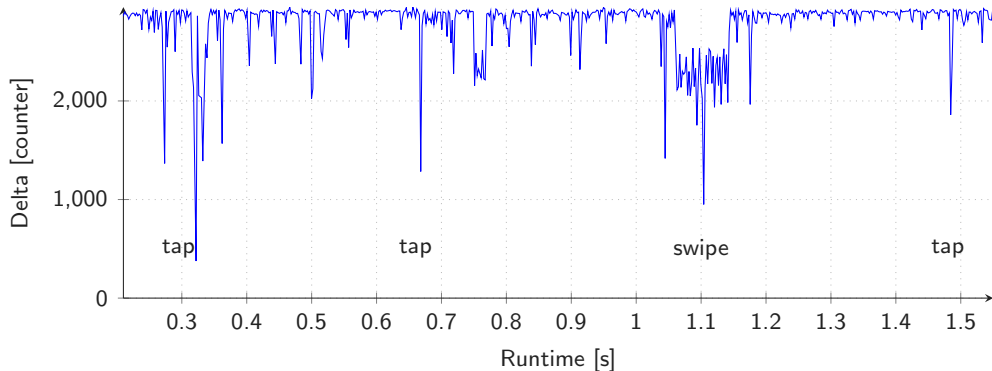


Figure 3: Keystroke timing on Google Nexus 5.

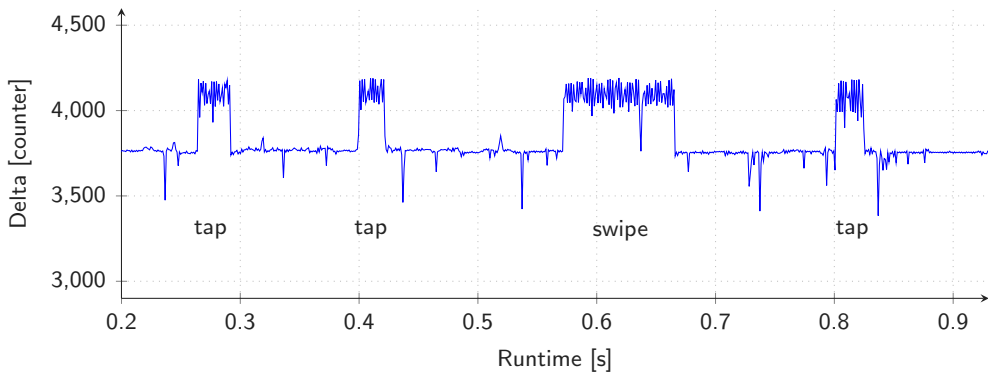


Figure 4: Keystroke timing on Xiaomi Redmi Note 3.

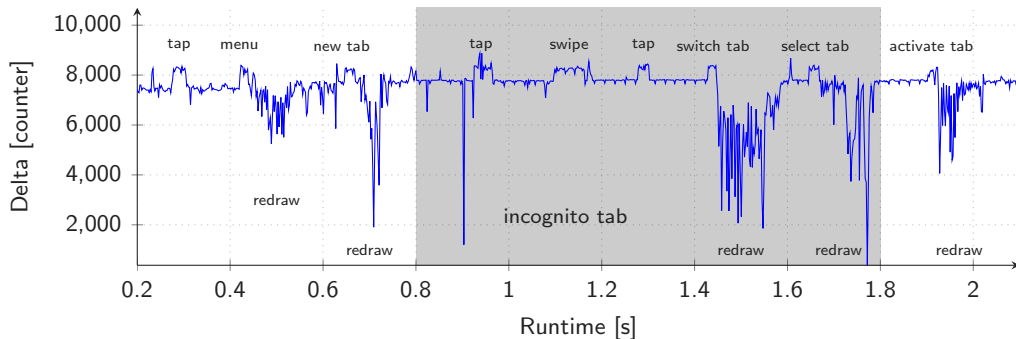


Figure 5: Chrome on Xiaomi Redmi Note 3.

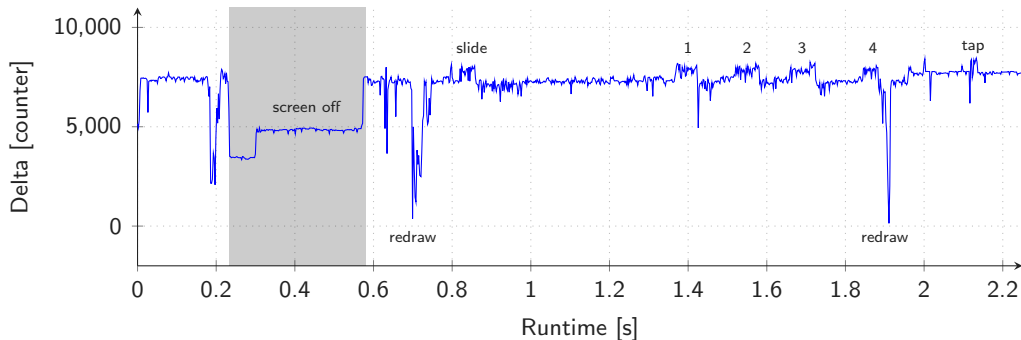


Figure 6: Firefox on Xiaomi Redmi Note 3.



- Timers were always the main focus



- Timers were always the main focus
- Reducing timer resolution is not sufficient



- Timers were always the main focus
- Reducing timer resolution is not sufficient
- Timers can (always) be built



- Timers were always the main focus
- Reducing timer resolution is not sufficient
- Timers can (always) be built
- Some attacks do not require timers at all



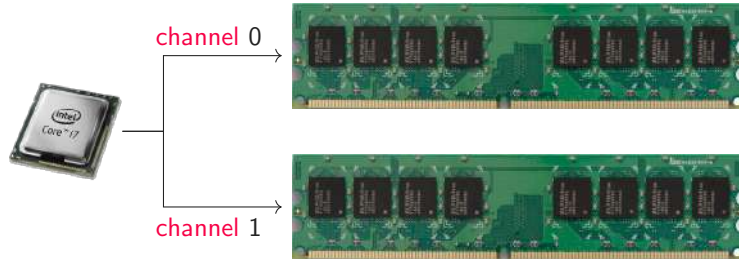
- Timers were always the main focus
- Reducing timer resolution is not sufficient
- Timers can (always) be built
- Some attacks do not require timers at all
- Important to understand requirements before designing countermeasures

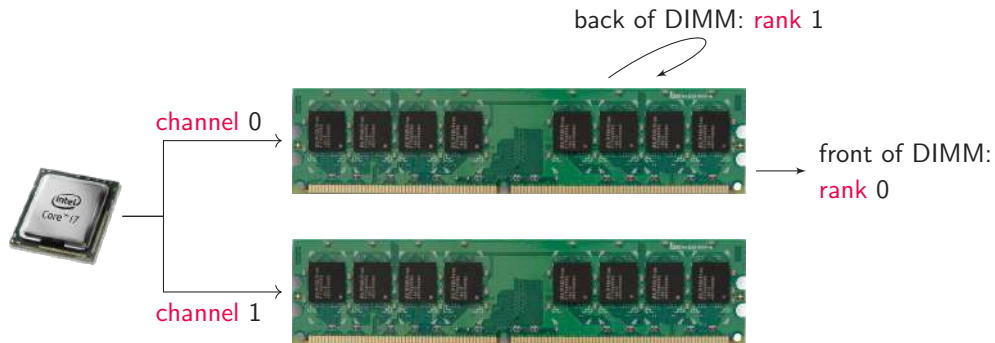
SIDE CHANNELS?

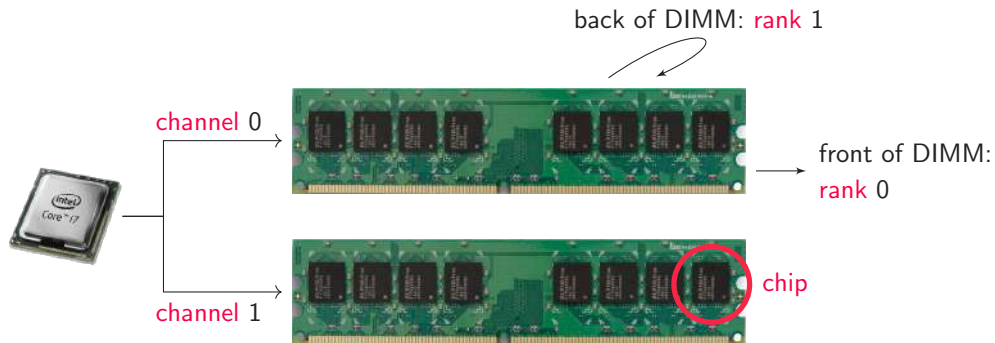


I WOULD NEED TO BREAK SOME ISOLATION

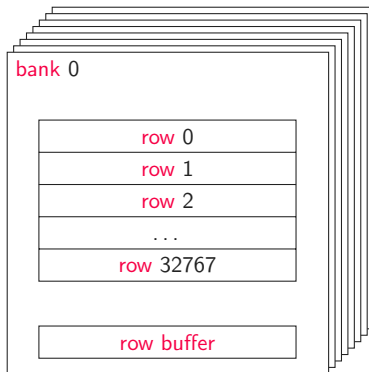




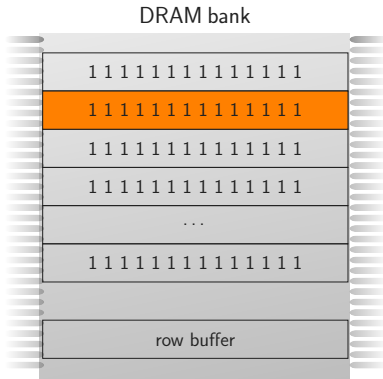


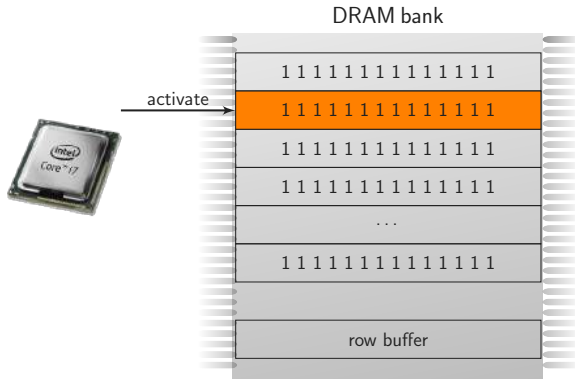


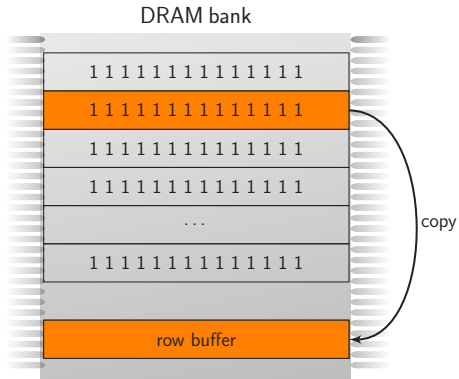
chip

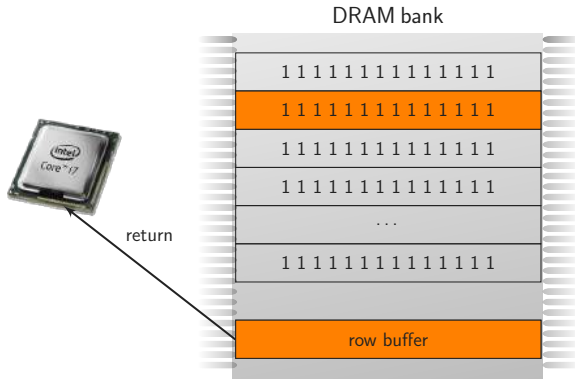


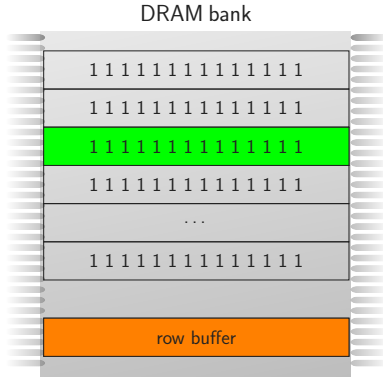
- bits in cells in rows
- access: **activate** row,
copy to row buffer

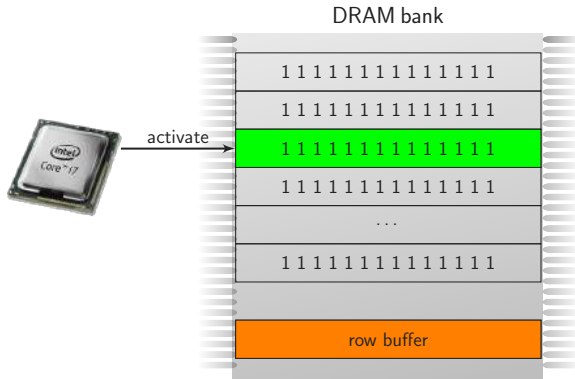


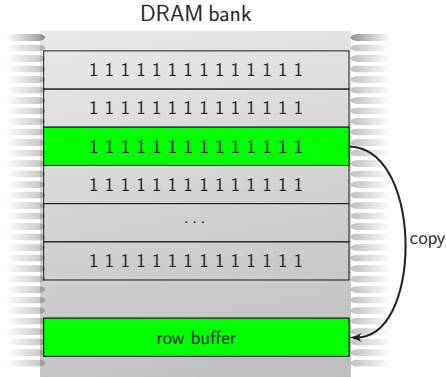


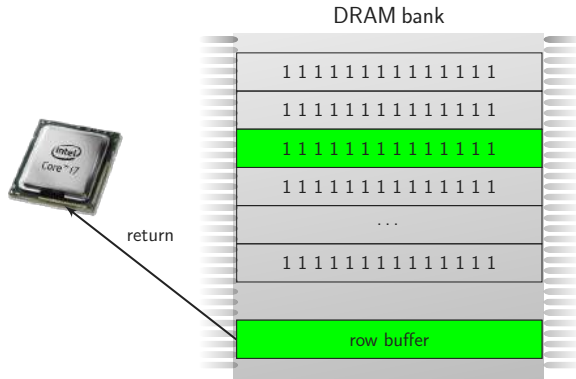


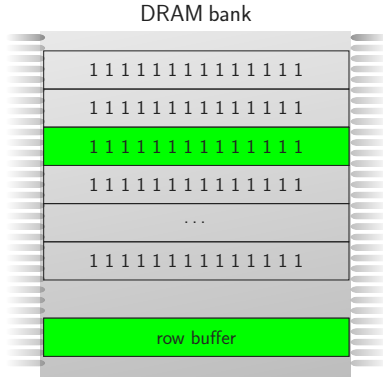


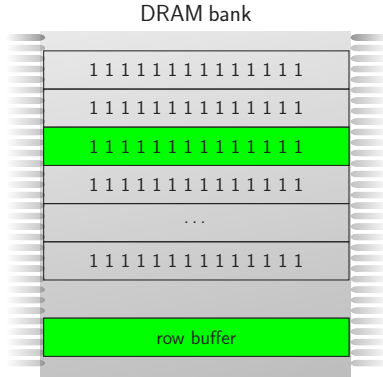


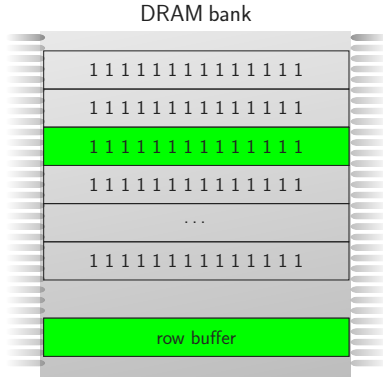


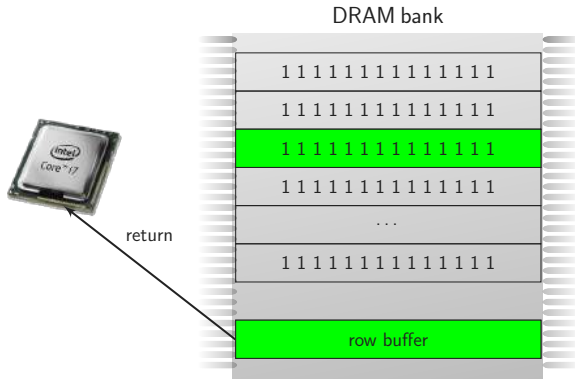


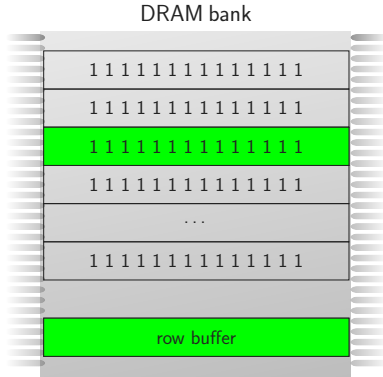












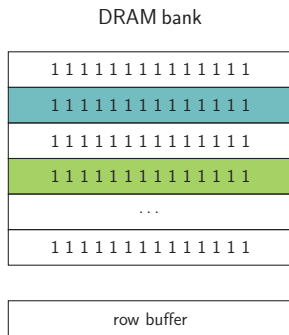
ONE DOES NOT SIMPLY



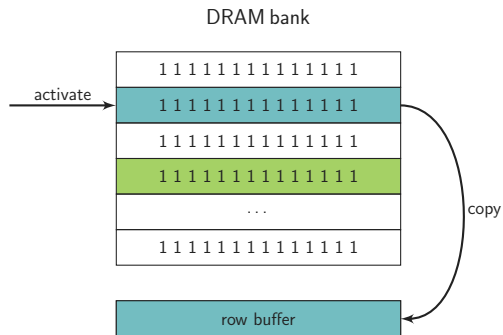
ROWHAMMER IN JAVASCRIPT

memegenerator.net

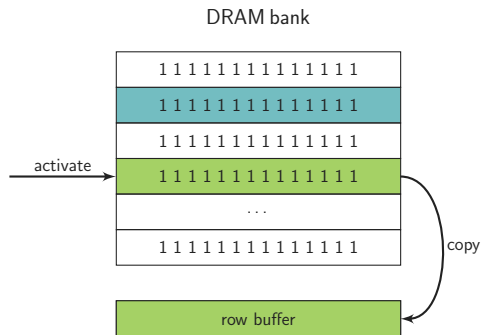
“It’s like breaking into an apartment by repeatedly slamming a neighbor’s door until the vibrations open the door you were after” – Motherboard Vice



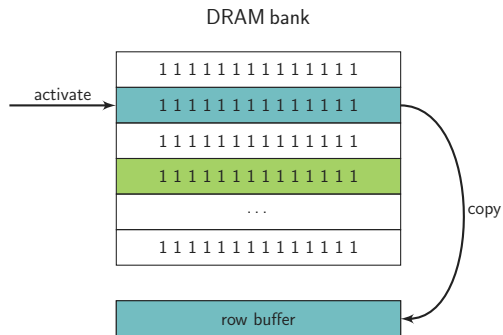
“It’s like breaking into an apartment by repeatedly slamming a neighbor’s door until the vibrations open the door you were after” – Motherboard Vice



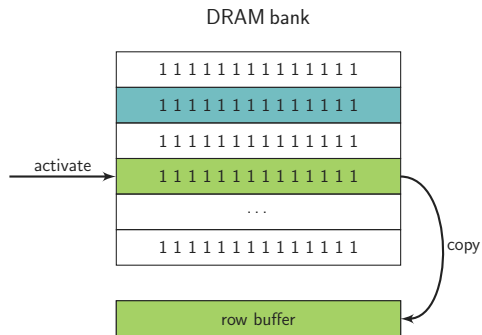
“It’s like breaking into an apartment by repeatedly slamming a neighbor’s door until the vibrations open the door you were after” – Motherboard Vice



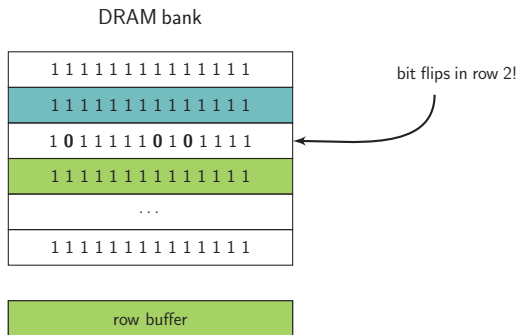
“It’s like breaking into an apartment by repeatedly slamming a neighbor’s door until the vibrations open the door you were after” – Motherboard Vice



“It’s like breaking into an apartment by repeatedly slamming a neighbor’s door until the vibrations open the door you were after” – Motherboard Vice



“It’s like breaking into an apartment by repeatedly slamming a neighbor’s door until the vibrations open the door you were after” – Motherboard Vice



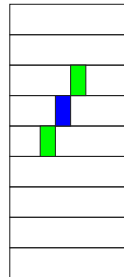
cache set 1

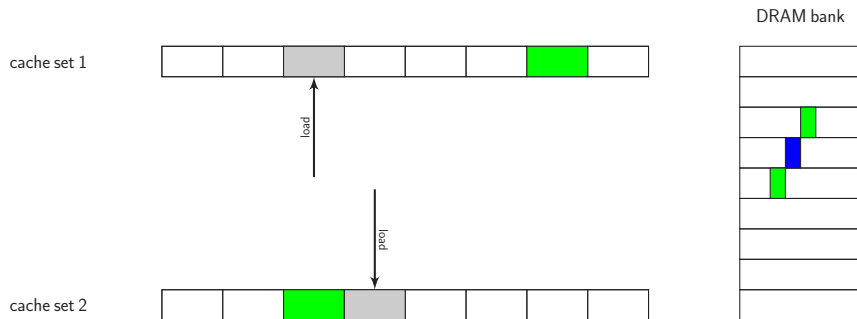


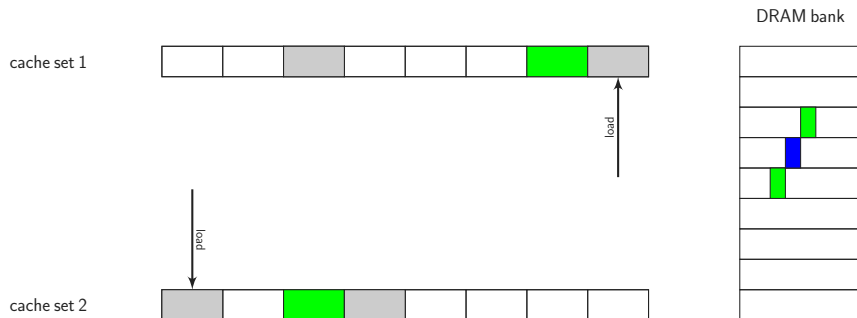
cache set 2

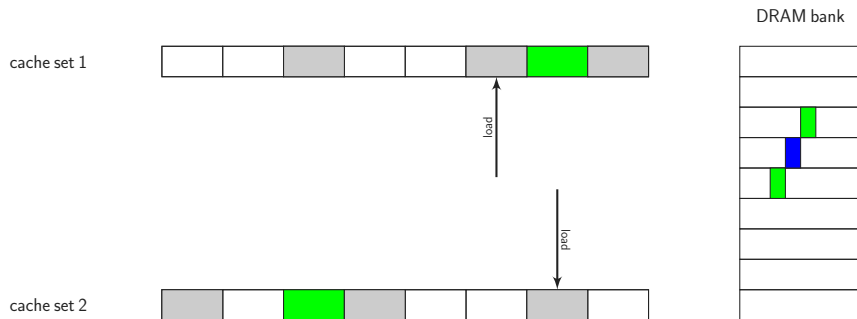


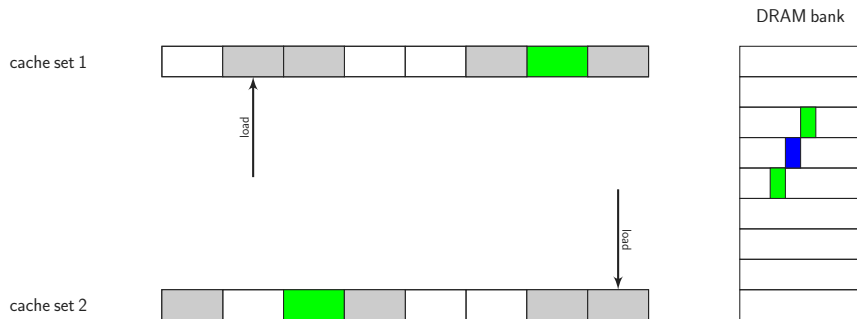
DRAM bank

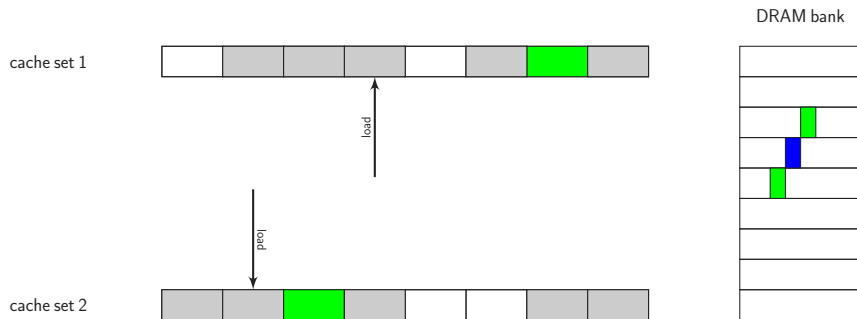


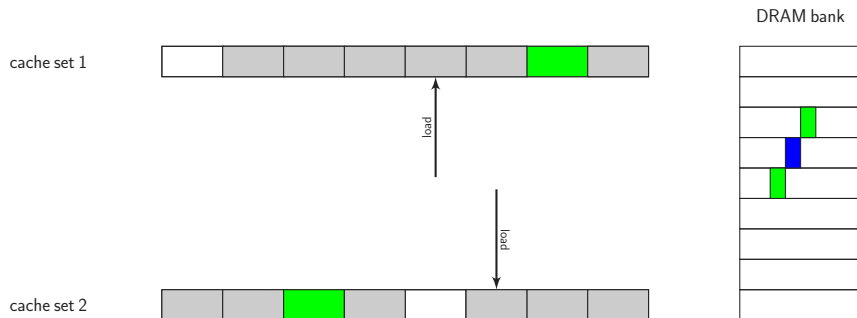


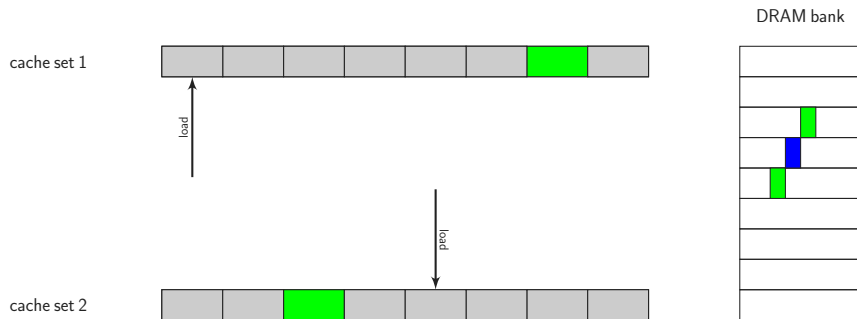


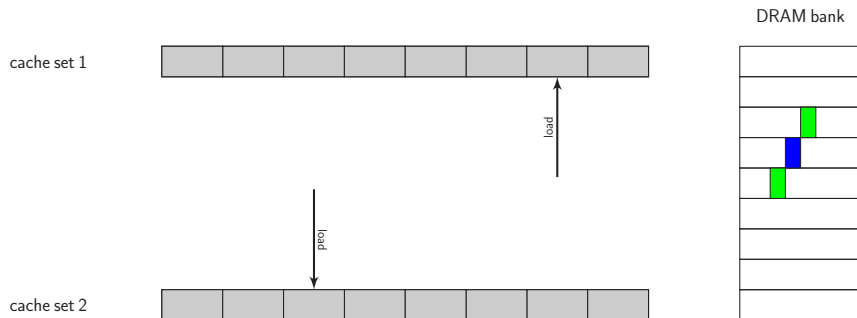


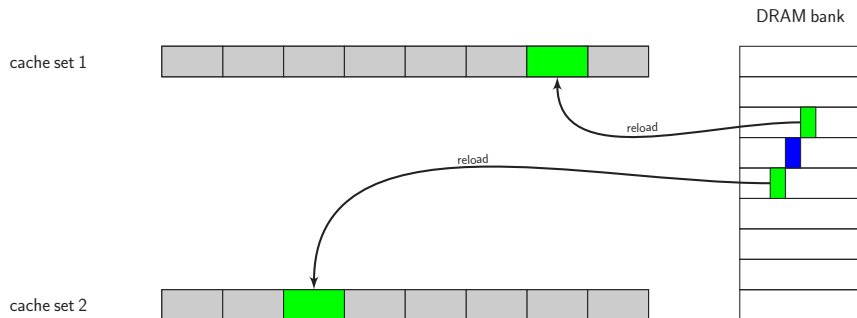


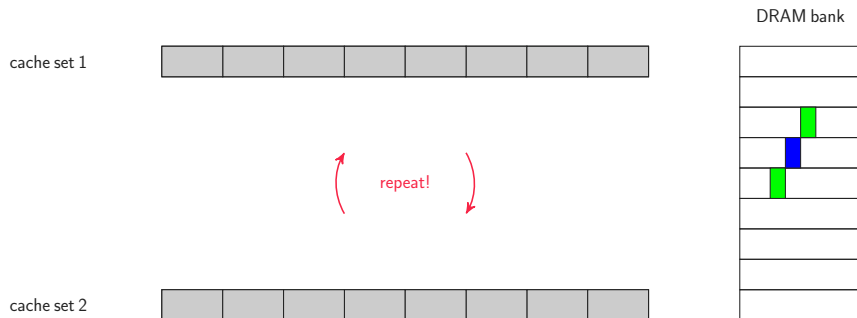


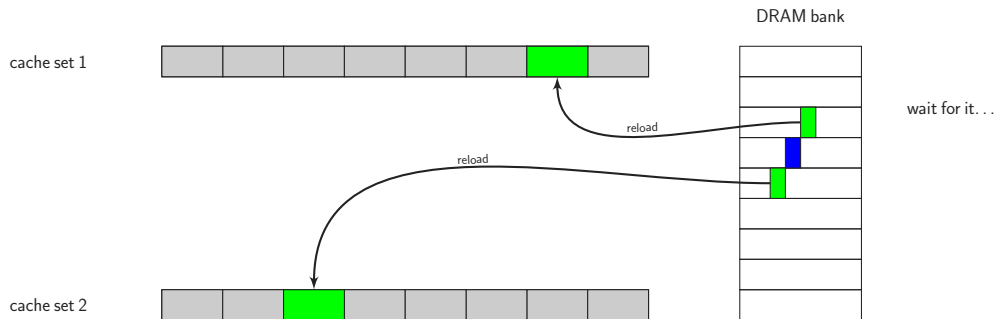












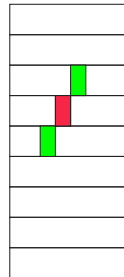
cache set 1



cache set 2



DRAM bank



bit flip!

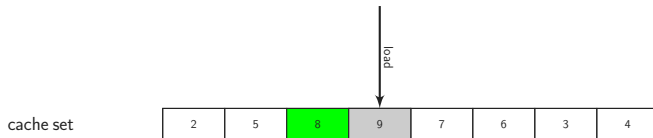
“LRU eviction” memory accesses: n accesses for an n -way cache

cache set

2	5	8	1	7	6	3	4
---	---	---	---	---	---	---	---

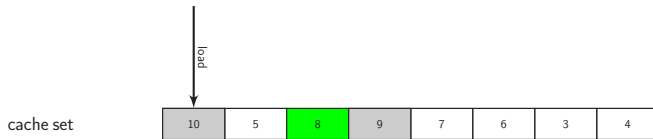
- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



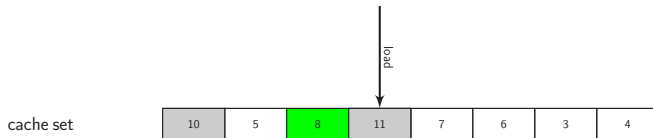
- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



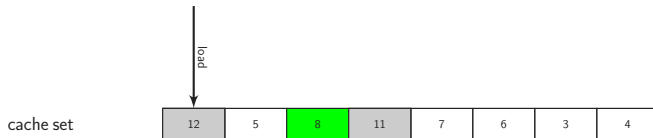
- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



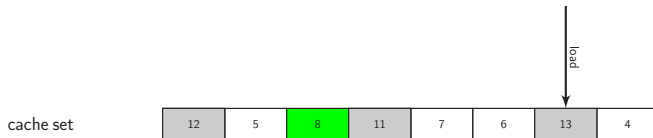
- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



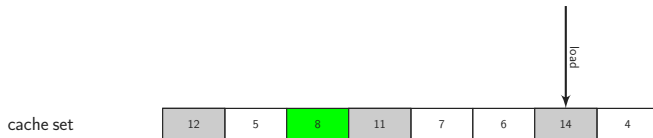
- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



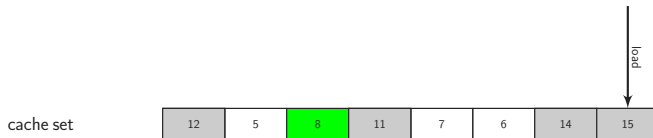
- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



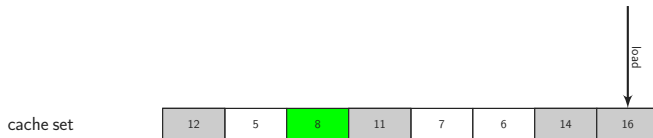
- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache



- no LRU replacement

“LRU eviction” memory accesses: n accesses for an n -way cache

cache set

12	5	8	11	7	6	14	16
----	---	---	----	---	---	----	----

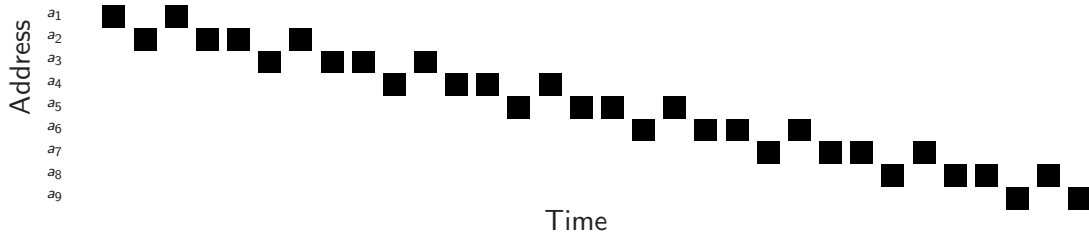
- no LRU replacement
- only 75% success rate on Haswell

“LRU eviction” memory accesses: n accesses for an n -way cache

cache set

12	5	8	11	7	6	14	16
----	---	---	----	---	---	----	----

- no LRU replacement
- only 75% success rate on Haswell
- more accesses → higher success rate, but **too slow**



→ fast and effective on Haswell: eviction rate $>99.97\%$

- represent accesses as a sequence: 1, 2, 1, 2, 2, 3, 2, 3, 3, 4, 3, 4, ...
- what can improve eviction rates?


- represent accesses as a sequence: 1, 2, 1, 2, 2, 3, 2, 3, 3, 4, 3, 4, ...
 - what can improve eviction rates?
- adding more *unique* addresses
- *more* accesses to the same addresses
- indistinguishable → *balanced* number of accesses

Write eviction strategies as: *P-C-D-L-S*

```
for (s = 0; s <= S - D; s += L)
  for (c = 0; c <= C; c += 1)
    for (d = 0; d <= D; d += 1)
      *a[s+d];
```

Write eviction strategies as: $P-C-D-L-S$

S : total number of different addresses
(= set size)




```
for (s = 0; s <= S - D; s += L)
  for (c = 0; c <= C; c += 1)
    for (d = 0; d <= D; d += 1)
      *a[s+d];
```

Write eviction strategies as: *P-C-D-L-S*

S: total number of different addresses
(= set size)

D: different addresses per inner access
loop



```
for (s = 0; s <= S - D; s += L)
  for (c = 0; c <= C; c += 1)
    for (d = 0; d <= D; d += 1)
      *a[s+d];
```

Write eviction strategies as: *P-C-D-L-S*

S: total number of different addresses
(= set size)

D: different addresses per inner access

```
for (s = 0; s <= S - D; s += L)
  for (c = 0; c <= C; c += 1)
    for (d = 0; d <= D; d += 1)
      *a[s+d];
```

Diagram illustrating the notation for cache eviction strategies using the example code snippet:

- S*: total number of different addresses (= set size) (points to the blue box *S* in the code)
- D*: different addresses per inner access (points to the orange box *D* in the code)
- C*: different addresses per inner access (points to the purple box *C* in the code)
- L*: step size of the inner access loop (points to the green box *L* in the code)

Write eviction strategies as: *P-C-D-L-S*

S: total number of different addresses
(= set size)

D: different addresses per inner access
loop

```
for (s = 0; s <= S - D; s += L)
  for (c = 0; c <= C; c += 1)
    for (d = 0; d <= D; d += 1)
      *a[s+d];
```

L: step size of the inner access
loop


C: number of repetitions of the inner
access loop


```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```

```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```

- $P-2-2-1-4 \rightarrow 1, 2, 1, 2, 2, 3, 2, 3, 3, 4, 3, 4$

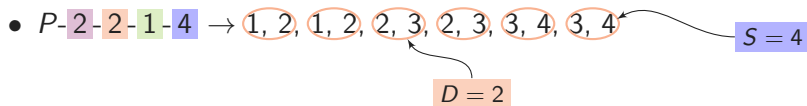
```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```

- $P-2-2-1-4 \rightarrow 1, 2, 1, 2, 2, 3, 2, 3, 3, 4, 3, 4$  $S = 4$

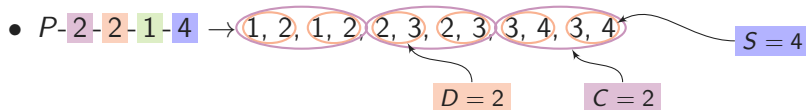
```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```

- $P-2-2-1-4 \rightarrow (1, 2), (1, 2), (2, 3), (2, 3), (3, 4), (3, 4)$ $\xleftarrow{S=4}$

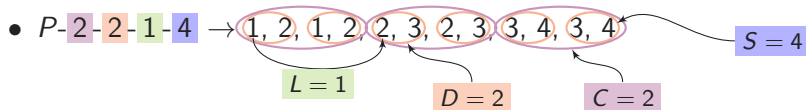
```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```



```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```



```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```



```
for (s = 0; s <= S - D; s += L)
  for (c = 1; c <= C; c += 1)
    for (d = 1; d <= D; d += 1)
      *a[s+d];
```

- $P-2-2-1-4 \rightarrow 1, 2, 1, 2, 2, 3, 2, 3, 3, 4, 3, 4$
- $P-1-1-1-4 \rightarrow 1, 2, 3, 4 \rightarrow$ LRU eviction with set size 4

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17		
<i>P</i> -1-1-1-20	20		

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	
<i>P</i> -1-1-1-20	20	99.82% ✓	

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗
<i>P</i> -2-1-1-17	34		

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗
<i>P</i> -2-1-1-17	34	99.86% ✓	

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗
<i>P</i> -2-1-1-17	34	99.86% ✓	191 ns ✓

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗
<i>P</i> -2-1-1-17	34	99.86% ✓	191 ns ✓
<i>P</i> -2-2-1-17	64		

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗
<i>P</i> -2-1-1-17	34	99.86% ✓	191 ns ✓
<i>P</i> -2-2-1-17	64	99.98% ✓	

¹Executed in a loop, on a Haswell with a 16-way last-level cache

We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗
<i>P</i> -2-1-1-17	34	99.86% ✓	191 ns ✓
<i>P</i> -2-2-1-17	64	99.98% ✓	180 ns ✓

¹Executed in a loop, on a Haswell with a 16-way last-level cache

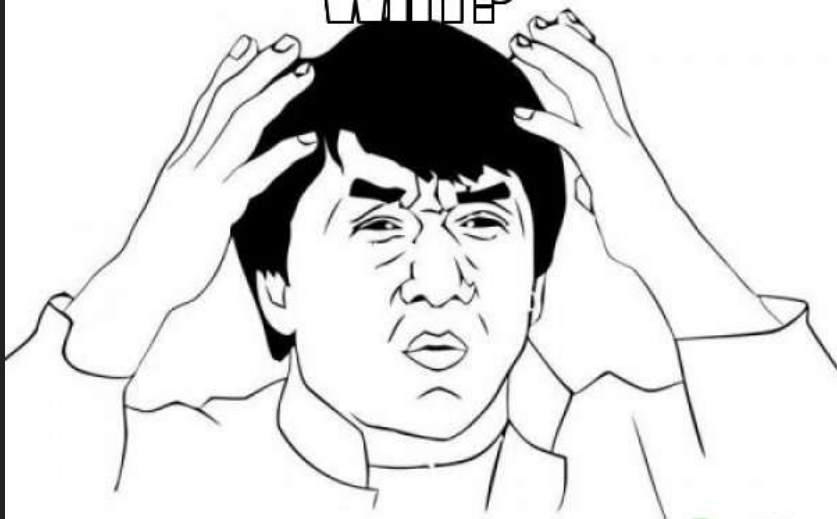
We evaluated more than 10000 strategies...¹

strategy	# accesses	eviction rate	loop time
<i>P</i> -1-1-1-17	17	74.46% ✗	307 ns ✓
<i>P</i> -1-1-1-20	20	99.82% ✓	934 ns ✗
<i>P</i> -2-1-1-17	34	99.86% ✓	191 ns ✓
<i>P</i> -2-2-1-17	64	99.98% ✓	180 ns ✓

→ more accesses, smaller execution time?

¹Executed in a loop, on a Haswell with a 16-way last-level cache

WHY?




Memes Happen

$P-1-1-1-17$ (17 accesses, 307ns)

$P-2-1-1-17$ (34 accesses, 191ns)

Time in ns



$P-1-1-1-17$ (17 accesses, 307ns)

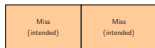


$P-2-1-1-17$ (34 accesses, 191ns)

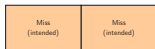


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

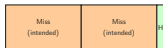


P -2-1-1-17 (34 accesses, 191ns)

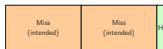


Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

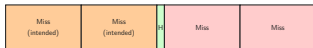


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

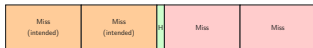


P -2-1-1-17 (34 accesses, 191ns)

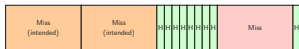


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

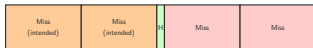


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

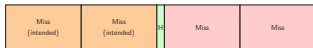


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)

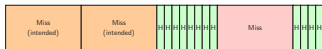


Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)

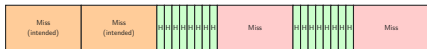


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

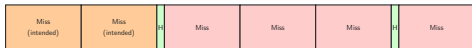


P -2-1-1-17 (34 accesses, 191ns)

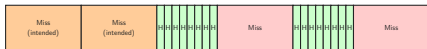


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

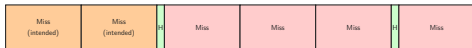


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

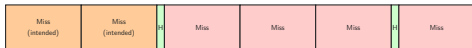


P -2-1-1-17 (34 accesses, 191ns)

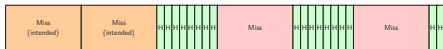


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

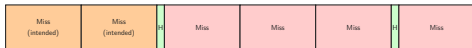


P -2-1-1-17 (34 accesses, 191ns)

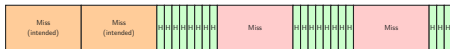


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

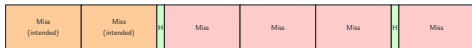


P -2-1-1-17 (34 accesses, 191ns)

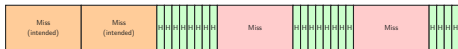


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

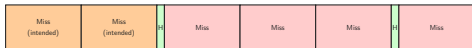


P -2-1-1-17 (34 accesses, 191ns)

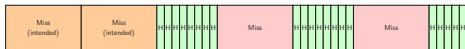


Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)

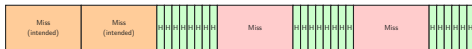


Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)

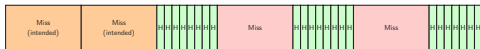


Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

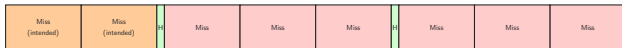


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

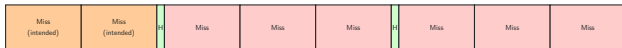


P -2-1-1-17 (34 accesses, 191ns)

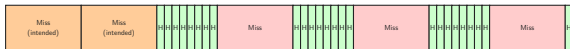


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

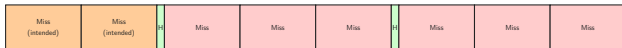


P -2-1-1-17 (34 accesses, 191ns)

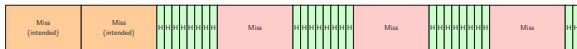


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

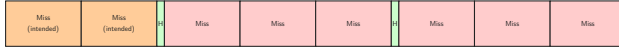


P -2-1-1-17 (34 accesses, 191ns)

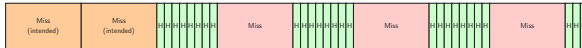


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

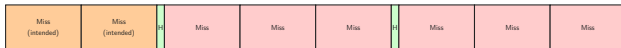


P -2-1-1-17 (34 accesses, 191ns)

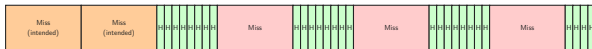


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

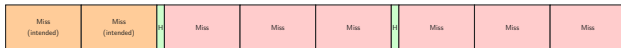


P -2-1-1-17 (34 accesses, 191ns)

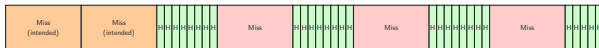


Time in ns

P -1-1-1-17 (17 accesses, 307ns)

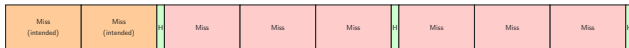


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

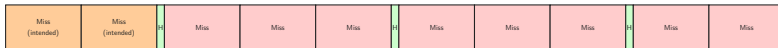


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)

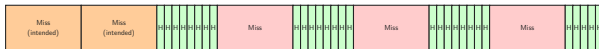


Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)

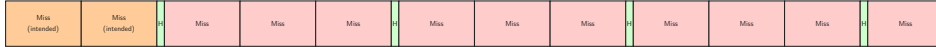


P -2-1-1-17 (34 accesses, 191ns)



Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)

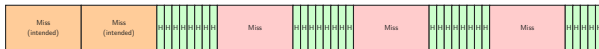


Time in ns

P -1-1-1-17 (17 accesses, 307ns)



P -2-1-1-17 (34 accesses, 191ns)



Time in ns

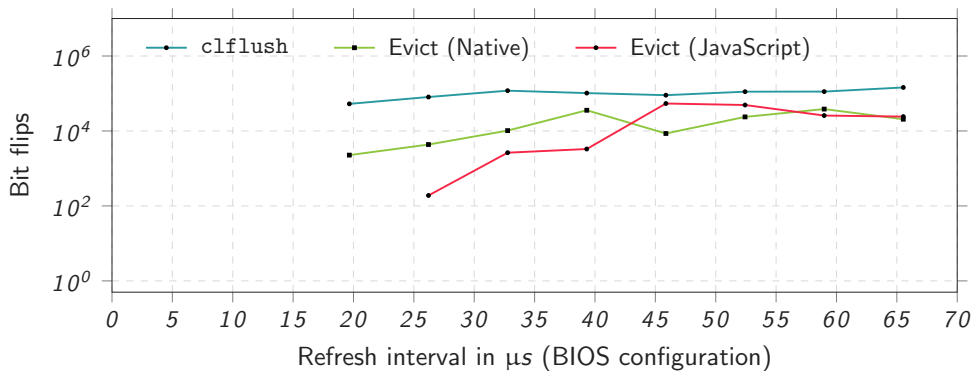


Figure 7: Number of bit flips within 15 minutes.



20: 12

30: 9

40: 1

50: 0

60: 1

70: 2

80: 199

90: 76

00: 72

10: 231

20: 572

250

] Found flip (254 != 255) at array index 340021386 when hammering indices 339881984 and 340156416

] Found flip (239 != 255) at array index 340022176 when hammering indices 339881984 and 340156416

] Found flip (191 != 255) at array index 340023138 when hammering indices 339881984 and 340156416

] Found flip (254 != 255) at array index 340025146 when hammering indices 339881984 and 340156416



ROOT privileges for web apps!

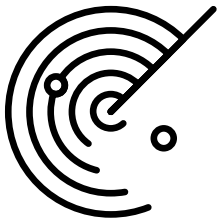
JavaScript zero



REAL
JavaScript
AND ZERO
SIDE-CHANNEL
ATTACKS

- Currently 11 microarchitectural and side-channel attacks in JavaScript





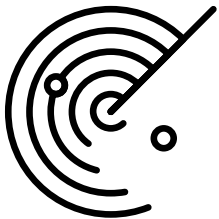
- Currently 11 microarchitectural and side-channel attacks in JavaScript
- Analyse requirements for every attack



- Currently **11** microarchitectural and side-channel **attacks** in JavaScript
- Analyse requirements for every attack
- Results in **5 categories**



- Currently **11** microarchitectural and side-channel **attacks** in JavaScript
- Analyse requirements for every attack
- Results in **5 categories**
 - Memory addresses
 - Accurate timing
 - Multithreading
 - Shared data
 - Sensor API



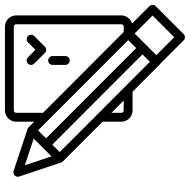
- Currently **11** microarchitectural and side-channel **attacks** in JavaScript
- Analyse requirements for every attack
- Results in **5 categories**
 - Memory addresses
 - Accurate timing
 - Multithreading
 - Shared data
 - Sensor API
- Every attack is in **at least one** category

	Mem. addrs.	Accurate timing	Multithreading	Shared data	Sensor API
Rowhammer.js	●	●	○	○	○
Practical Memory Deduplication Attacks in Sandboxed Javascript	◐	●	○	○	○
Fantastic Timers and Where to Find Them	●	● [†]	◐	◐	○
ASLR on the Line	●	● [†]	◐	◐	○
The spy in the sandbox	◐	●	○	○	○
Loophole	○	◐	●	○	○
Pixel perfect timing attacks with HTML5	○	● [†]	◐	◐	○
The clock is still ticking	○	●	◐	○	○
Practical Keystroke Timing Attacks in Sandboxed JavaScript	○	◐ [†]	●	◐	○
TouchSignatures	○	○	○	○	●
Stealing sensitive browser data with the W3C Ambient Light Sensor API	○	○	○	○	●

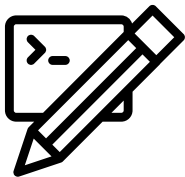
[†] If accurate timing is not available, it can be approximated using a combination of multithreading and shared

Prevents Defense	Rowham- mer.js	Page Dedu- plication	DRAM Covert Channel	Anti- ASLR	Cache Eviction	Keystroke Timing	Browser
Buffer ASLR	○	◐	○	●	●	○	○
Array preloading	●	○	●	○	○	○	○
Non-deterministic array	●	◐	◐	●	●	○	○
Array index randomization	○	●	○	●	○	○	○
Low-resolution timestamp	○	◐	○	○	○	◐	◐
Fuzzy time	○	◐*	○	○*	○	●*	●*
WebWorker polyfill	○	○	●	●	●	●	○
Message delay	○	○	○	○	○	◐	◐
Slow SharedArrayBuffer	○	○	●	◐	●	○	○
No SharedArrayBuffer	○	○*	●	●*	●	○*	○*
Summary	●	●	●	●	●	●	●

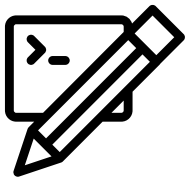
Prevented (●), partly prevented / more difficult (◐), not prevented (○). A star (*) means the combination is necessary.



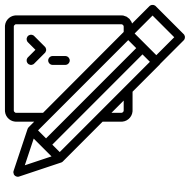
- Ideally → browser core



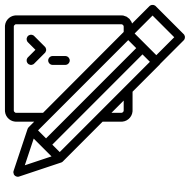
- Ideally → browser core
- Maintaining a fork is hard



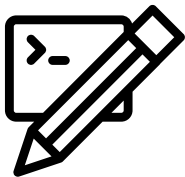
- Ideally → browser core
- Maintaining a fork is hard
- Generic solution for multiple browsers



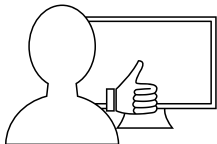
- Ideally → **browser core**
- Maintaining a fork is hard
- Generic solution for multiple browsers
- Parsing JavaScript is hard



- Ideally → browser core
- Maintaining a fork is hard
- Generic solution for multiple browsers
- Parsing JavaScript is hard
- Implementation in JavaScript → Virtual machine layering



- Ideally → **browser core**
- Maintaining a fork is hard
- Generic solution for multiple browsers
- Parsing JavaScript is hard
- Implementation in JavaScript → **Virtual machine layering**
- Proof-of-concept → browser extension



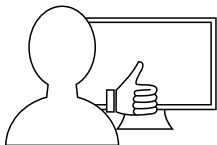
- Affects user experience? E.g., disable multithreading



- Affects user experience? E.g., disable multithreading
- Select pre-defined **protection level**

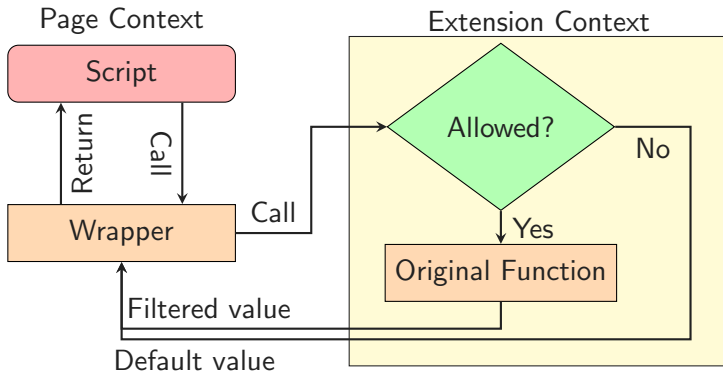


- Affects user experience? E.g., disable multithreading
- Select pre-defined **protection level**
- Protection levels → combinations of defenses



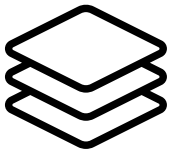
- Affects user experience? E.g., disable multithreading
- Select pre-defined **protection level**
- Protection levels → combinations of defenses
- Each defense is disabled, enabled, or prompts

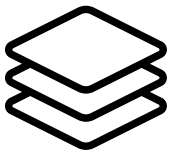
- Functions and properties are replaced by **wrappers**



- Functions can be **re-defined** in JavaScript

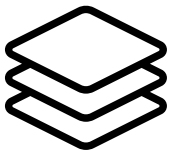
```
1 var original_reference = window.performance.now;  
2 window.performance.now = function () { return 0; };  
3
```





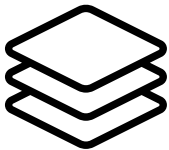
- Functions can be **re-defined** in JavaScript

```
1 var original_reference = window.performance.now;  
2 window.performance.now = function() { return 0; };  
3  
4 // call the new function (via function name)  
5 alert(window.performance.now()); // == alert(0)  
6
```



- Functions can be **re-defined** in JavaScript

```
1 var original_reference = window.performance.now;  
2 window.performance.now = function() { return 0; };  
3  
4 // call the new function (via function name)  
5 alert(window.performance.now()); // == alert(0)  
6  
7 // call the original function (only via reference)  
8 alert(original_reference.call(window.performance));
```

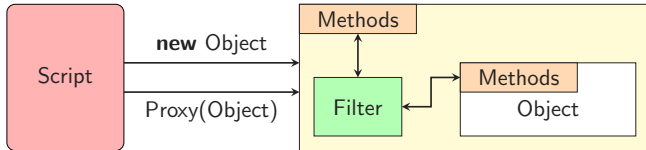


- Functions can be **re-defined** in JavaScript

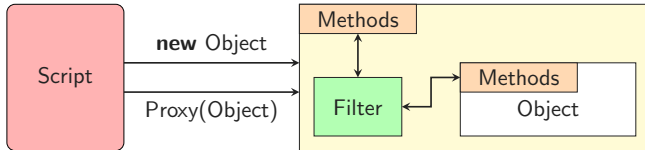
```
1 var original_reference = window.performance.now;  
2 window.performance.now = function() { return 0; };  
3  
4 // call the new function (via function name)  
5 alert(window.performance.now()); // == alert(0)  
6  
7 // call the original function (only via reference)  
8 alert(original_reference.call(window.performance));
```

- Properties can be replaced by **accessor properties**

- Objects are proxied

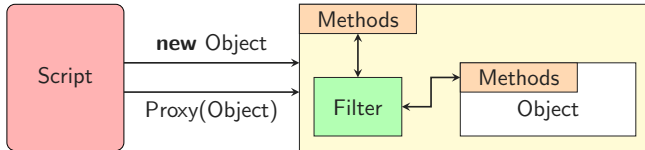


- Objects are proxied



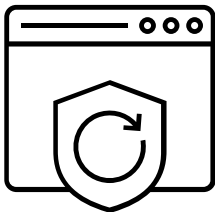
- All properties and functions are handled by the original object

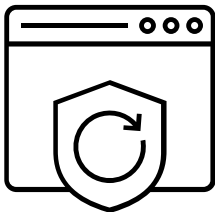
- Objects are proxied



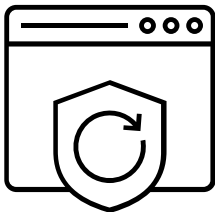
- All properties and functions are handled by the original object
- Functions and properties can be overwritten in the **proxy object**

- Attacker tries to circumvent JavaScript Zero



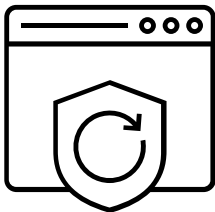


- Attacker tries to circumvent JavaScript Zero
- Self protection is necessary if implemented in JavaScript



- Attacker tries to circumvent JavaScript Zero
- **Self protection** is necessary if implemented in JavaScript
- Use closures to hide all references to original functions

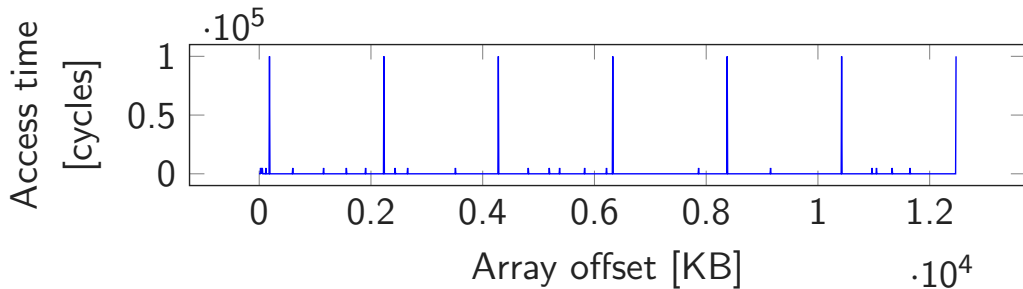
```
1 (function() {  
2 // original is only accessible in this scope  
3 var original = window.performance.now;  
4 window.performance.now = ...  
5 })();
```

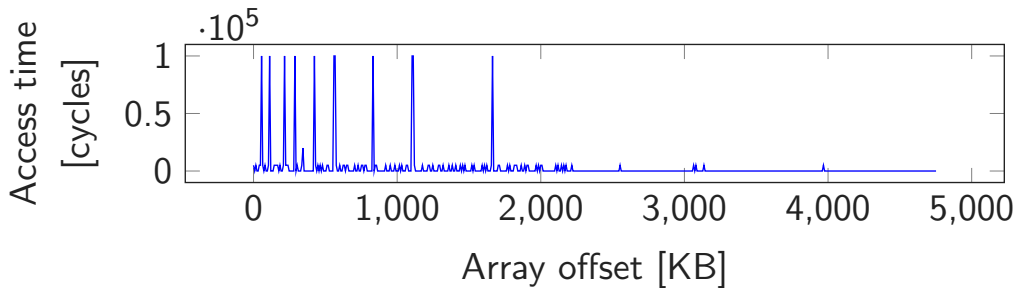


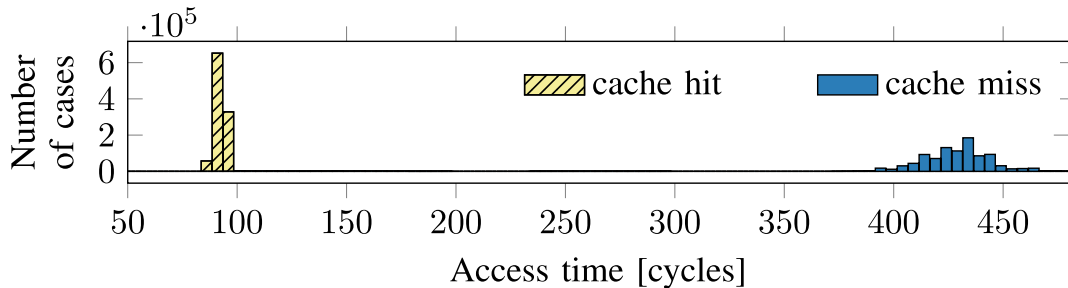
- Attacker tries to circumvent JavaScript Zero
- **Self protection** is necessary if implemented in JavaScript
- Use closures to hide all references to original functions

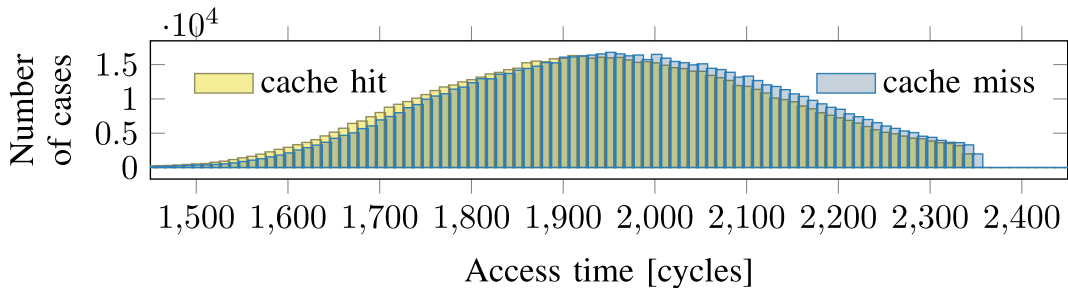
```
1 (function() {  
2 // original is only accessible in this scope  
3 var original = window.performance.now;  
4 window.performance.now = ...  
5 })();
```

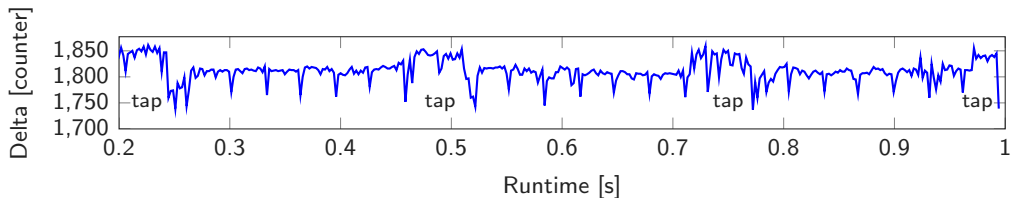
- Prevent objects from being modified: `Object.freeze`

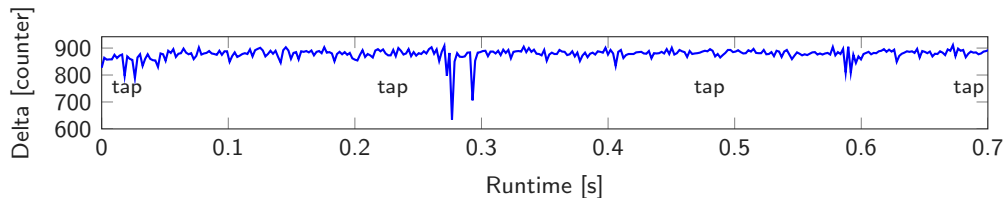


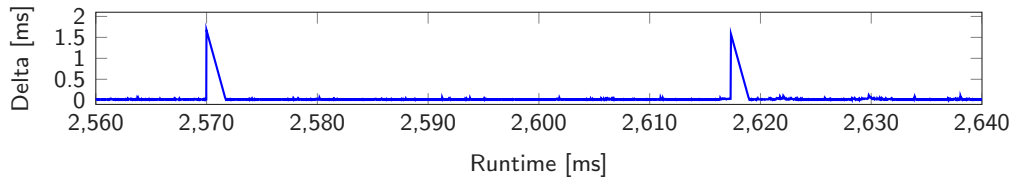


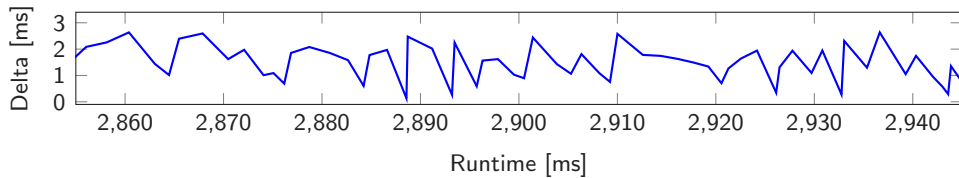


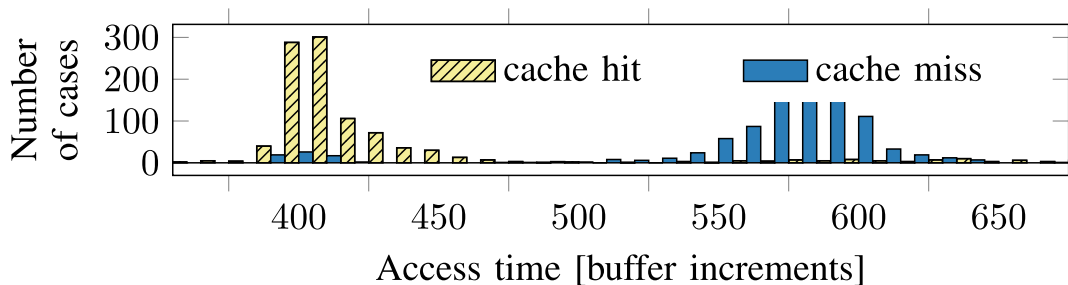


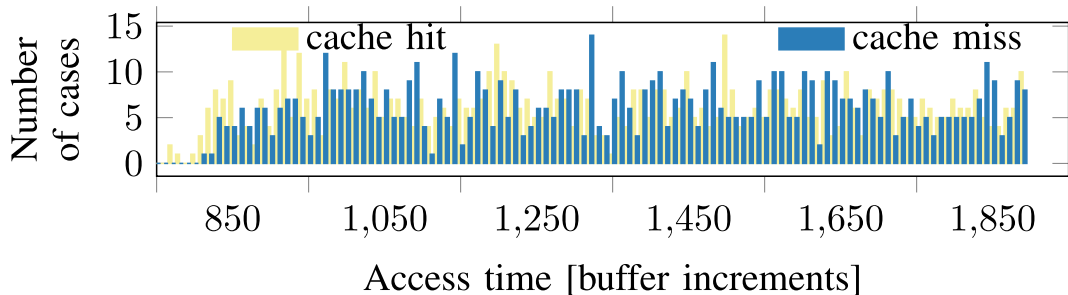


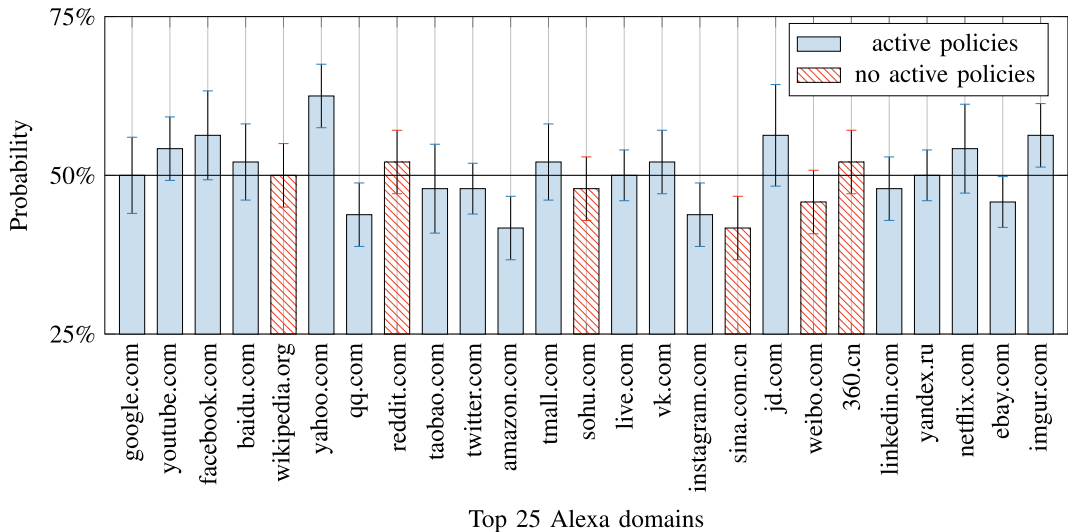




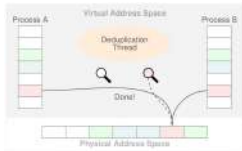


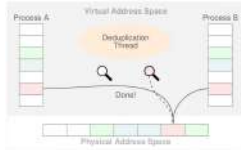


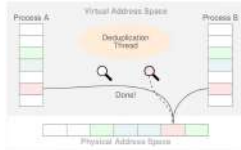


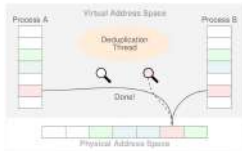


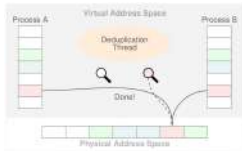


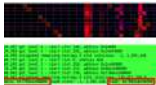
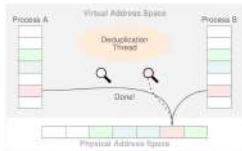


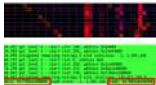
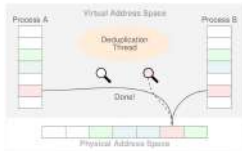












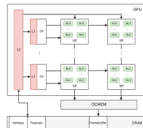
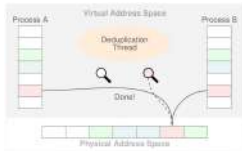


Fig. 2: Building blocks of an integrated GPU

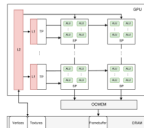
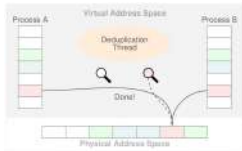


Fig. 2: Building blocks of an integrated GPU



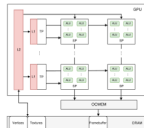
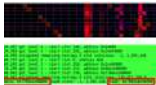
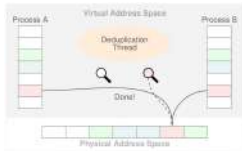


Fig. 2: Building blocks of an integrated GPU





- Jumping Abstraction Layers becomes easier



- Jumping Abstraction Layers becomes easier
- New attacks often also in JavaScript



- Jumping Abstraction Layers becomes easier
- New attacks often also in JavaScript
- We need cross-layer solutions

Jumping Abstraction Layers: Microarchitectural Attacks in JavaScript

Daniel Gruss

September 18, 2019

Graz University of Technology